

Circuit ORAM

Xiao Shaun Wang (UMD), T-H. Hubert Chan (HKU),
and Elaine Shi (UMD)





ORAM in
Secure
Computation

US Government Investment in MPC:

NSF: ~\$25M

DARPA: ~\$25M

AFOSR: ~\$15M

IARPA, NSA: ? M

[Gordon et al. 13], [Gentry et al. 13], [Liu et al. 13],
[Gentry et al. 14], [Keller et al. 14], [Wang et al. 14], etc.



Background

ORAM has been
optimized for a
wrong metric.

w.r.t. **secure computation**

Traditional metric: bandwidth overhead

Metric for secure computation: **Circuit Size**

ORAMs with
Small Circuit
Complexity?

Circuit ORAM achieves
 $O(D \log N)$ circuit
complexity for blocks
of size

$D = \Omega(\log^2 N)$ bits

Smallest circuit size both **asymptotically** and
in practice.



Empirical
Results

Circuit ORAM
outperforms Path
ORAM by **8x - 48x** at
1 GB data size.

Speedup depends on what variations of Path
ORAM is used.



Our Hope

ORAM accesses may be
securely evaluated
potentially at **hundreds**
of accesses/sec
for **4 MB** data size
(assuming certain offline preparation)

Garbling can be done at **10^8 gates/sec** using
off-the-shelf modern processors

(not counting other overhead such as OT)

[Bellare et al. 13]

On the Tightness
of the ORAM
Lower Bound

Circuit ORAM:

For any $0 < \epsilon < 1$, any N -word RAM program with block size of $\Omega(N^\epsilon)$ can be simulated obliviously with $O(\log N)$ runtime blowup, with **inverse poly** failure probability.

[Goldreich 87, **stronger** interpretation]: $\Omega(\log N)$ runtime blowup is necessary for **any block size** and tolerate up to **constant** failure probability.

Consider $O(1)$ client storage

We are currently implementing Circuit ORAM
over garbled circuits!

Thank You

wangxiao@cs.umd.edu